# **6** Security & Monitorin<u>q</u>

In a traditional wired network, access control is very straightforward: If a person has physical access to a computer or network hub, they can use (or abuse) the network resources. While software mechanisms are an important component of network security, limiting physical access to the network devices is the ultimate access control mechanism. Simply put, if all terminals and network components are only accessible to trusted individuals, the network can likely be trusted.

The rules change significantly with wireless networks. While the apparent range of your access point may seem to be just a few hundred meters, a user with a high gain antenna may be able to make use of the network from several blocks away. Should an unauthorized user be detected, is impossible to simply "trace the cable" back to the user's location. Without transmitting a single packet, a nefarious user can even log all network data to disk. This data can later be used to launch a more sophisticated attack against the network. Never assume that radio waves simply "stop" at the edge of your property line.

It is usually unreasonable to completely trust all users of the network, even on wired networks. Disgruntled employees, uneducated network users, and simple mistakes on the part of honest users can cause significant harm to network operations. As the network architect, your goal is to facilitate private communication between legitimate users of the network. While a certain amount of access control and authentication is necessary in any network, you have failed in your job if legitimate users find it difficult to use the network to communicate.

There's an old saying that the only way to completely secure a computer is to unplug it, lock it in a safe, destroy the key, and bury the whole thing in concrete. While such a system might be completely "secure", it is useless for communication. When you make security decisions for your network, remember that above all else, the network exists so that its users can communicate with each other. Security considerations are important, but should not get in the way of the network's users.

# Physical security

When installing a network, you are building an infrastructure that people depend on. Security measures exist to ensure that the network is reliable. For many installations, outages often occur due to human tampering, whether accidental or not. Networks have physical components, such as wires and boxes, which are easily disturbed. In many installations, people will not understand the purpose of the installed equipment, or curiosity may lead them to experiment. They may not realize the importance of a cable connected to a port. Someone may unplug an Ethernet cable so that they can connect their laptop for 5 minutes, or move a switch because it is in their way. A plug might be removed from a power bar because someone needs that receptacle. Assuring the physical security of an installation is paramount. Signs and labels will only be useful to those who can read your language. Putting things out of the way and limiting access is the best means to assure that accidents and tinkering do not occur.

In less developed economies, proper fasteners, ties, or boxes will not be as easy to find. You should be able to find electrical supplies that will work just as well. Custom enclosures are also easy to manufacture and should be considered essential to any installation. It is often economical to pay a mason to make holes and install conduit. Where this would be an expensive option in the developed world, this type of labour intensive activity can be affordable in Southern countries. PVC can be embedded in cement walls for passing cable from room to room. This avoids the need to smash new holes every time a cable needs to be passed. Plastic bags can be stuffed into the conduit around the cables for insulation.

Small equipment should be mounted on the wall and larger equipment should be put in a closet or in a cabinet.

# Switches

Switches, hubs or interior access points can be screwed directly onto a wall with a wall plug. It is best to put this equipment as high as possible to reduce the chance that someone will touch the device or its cables.

# Cables

At the very least, cables should be hidden and fastened. It is possible to find plastic cable conduit that can be used in buildings. If you cannot find it, simple cable attachments can be nailed into the wall to secure the cable. This will make sure that the cable doesn't hang where it can be snagged, pinched or cut.

It is preferable to bury cables, rather than to leave them hanging across a yard. Hanging wires might be used for drying clothes, or be snagged by a ladder, etc. To avoid vermin and insects, use plastic electrical conduit. The marginal expense will be well worth the trouble. The conduit should be buried about 30 cm deep, or below the frost level in cold climates. It is worth the extra investment of buying larger conduit than is presently required, so that future cables can be run through the same tubing. Consider labeling buried cable with a "call before you dig" sign to avoid future accidental outages.

# Power

It is best to have power bars locked in a cabinet. If that is not possible, mount the power bar under a desk, or on the wall and use duct tape (or gaffer tape, a strong adhesive tape) to secure the plug into the receptacle. On the UPS and power bar, do not leave any empty receptacles. Tape them if necessary. People will have the tendency to use the easiest receptacle, so make these critical ones difficult to use. If you do not, you might find a fan or light plugged into your UPS; though it is nice to have light, it is nicer to keep your server running!

# Water

Protect your equipment from water and moisture. In all cases make sure that your equipment, including your UPS is at least 30 cm from the ground, to avoid damage from flooding. Also try to have a roof over your equipment, so that water and moisture will not fall onto it. In moist climates, it is important that the equipment has proper ventilation to assure that moisture can be exhausted. Small closets need to have ventilation, or moisture and heat can degrade or destroy your gear.

# Masts

Equipment installed on a mast is often safe from thieves. Nevertheless, to deter thieves and to keep your equipment safe from winds it is good to overengineer mounts. Painting equipment a dull white or grey color reflects the sun and makes it look plain and uninteresting. Panel antennas are often preferred because they are much more subtle and less interesting than dishes. Any installation on walls should be high enough to require a ladder to reach. Try choosing well-lit but not prominent places to put equipment. Also avoid antennae that resemble television antennae, as those are items that will attract interest by thieves, where a wifi antenna will be useless to the average thief.

# Threats to the network

One critical difference between Ethernet and wireless is that wireless networks are built on a *shared medium*. They more closely resemble the old network hubs than modern switches, in that every computer connected to the network can "see" the traffic of every other user. To monitor all network traffic on an access point, one can simply tune to the channel being used, put the network card into monitor mode, and log every frame. This data might be directly valuable to an eavesdropper (including data such as email, voice data, or online chat logs). It may also provide passwords and other sensitive data, making it possible to compromise the network even further. As we'll see later in this chapter, this problem can be mitigated by the use of encryption.

Another serious problem with wireless networks is that its users are relatively *anonymous*. While it is true that every wireless device includes a unique MAC address that is supplied by the manufacturer, these addresses can often be changed with software. Even when the MAC address is known, it can be very difficult to judge where a wireless user is physically located. Multipath effects, high-gain antennas, and widely varying radio transmitter characteristics can make it impossible to determine if a malicious wireless user is sitting in the next room or is in an apartment building a mile away.

While unlicensed spectrum provides a huge cost savings to the user, it has the unfortunate side effect that **denial of service** (**DoS**) attacks are trivially simple. By simply turning on a high powered access point, cordless phone, video transmitter, or other 2.4 GHz device, a malicious person could cause significant problems on the network. Many network devices are vulnerable to other forms of denial of service attacks as well, such as disassociation flooding and ARP table overflows.

Here are several categories of individuals who may cause problems on a wireless network:

• Unintentional users. As more wireless networks are installed in densely populated areas, it is common for laptop users to accidentally associate to the wrong network. Most wireless clients will simply choose any available wireless network when their preferred network is unavailable. The user may then make use of this network as usual, completely unaware that they may be transmitting sensitive data on someone else's network. Malicious people may even take advantage of this by setting up access points in strategic locations, to try to attract unwitting users and capture their data.

The first step in avoiding this problem is educating your users, and stressing the importance of connecting only to known and trusted networks. Many wireless clients can be configured to only connect to trusted networks, or to ask permission before joining a new network. As we will see later in this chapter, users can safely connect to open public networks by using strong encryption.

• *War drivers*. The "war driving" phenomenon draws its name from the popular 1983 hacker film, "War Games". War drivers are interested in finding the physical location of wireless networks. They typically drive around with a laptop, GPS, and omnidirectional antenna, logging the name and location of any networks they find. These logs are then combined with logs from other war drivers, and are turned into graphical maps depicting the wireless "footprint" of a particular city.

The vast majority of war drivers likely pose no direct threat to networks, but the data they collect might be of interest to a network cracker. For example, it might be obvious that an unprotected access point detected by a war driver is located inside a sensitive building, such as a government or corporate office. A malicious person could use this information to illegally access the network there. Arguably, such an AP should never have been set up in the first place, but war driving makes the problem all the more urgent. As we will see later in this chapter, war drivers who use the popular program NetStumbler can be detected with programs such as Kismet. For more information about war driving, see sites such as *http://www.wifimaps.com/*, *http://www.nodedb.com/*, or *http://www.netstumbler.com/*.

• **Rogue access points**. There are two general classes of rogue access points: those incorrectly installed by legitimate users, and those installed by malicious people who intend to collect data or do harm to the network. In the simplest case, a legitimate network user may want better wireless coverage in their office, or they might find security restrictions on the corporate wireless network too difficult to comply with. By installing an inexpensive consumer access point without permission, the user opens the entire network up to potential attacks from the inside. While it is possible to scan for unauthorized access points on your wired network, setting a clear policy that prohibits them is very important.

The second class of rogue access point can be very difficult to deal with. By installing a high powered AP that uses the same ESSID as an existing network, a malicious person can trick people into using their equipment, and log or even manipulate all data that passes through it. Again, if your users are trained to use strong encryption, this problem is significantly reduced.

• *Eavesdroppers*. As mentioned earlier, eavesdropping is a very difficult problem to deal with on wireless networks. By using a passive monitoring tool (such as Kismet), an eavesdropper can log all network data from a great distance away, without ever making their presence known. Poorly

encrypted data can simply be logged and cracked later, while unencrypted data can be easily read in real time.

If you have difficulty convincing others of this problem, you might want to demonstrate tools such as Etherpeg (*http://www.etherpeg.org/*) or Driftnet (*http://www.ex-parrot.com/~chris/driftnet/*). These tools watch a wireless network for graphical data, such as GIF and JPEG files. While other users are browsing the Internet, these tools simply display all graphics found in a graphical collage. I often use tools such as this as a demonstration when lecturing on wireless security. While you can tell a user that their email is vulnerable without encryption, nothing drives the message home like showing them the pictures they are looking at in their web browser.

Again, while it cannot be completely prevented, proper application of strong encryption will discourage eavesdropping.

This introduction is intended to give you an idea of the problems you are up against when designing a wireless network. Later in this chapter, we will look at tools and techniques that will help you to mitigate these problems.

# Authentication

Before being granted access to network resources, users should first be *authenticated*. In an ideal world, every wireless user would have an identifier that is unique, unchangeable, and cannot be impersonated by other users. This turns out to be a very difficult problem to solve in the real world.

The closest feature we have to a unique identifier is the MAC address. This is the 48-bit number assigned by the manufacturer to every wireless and Ethernet device. By employing *mac filtering* on our access points, we can authenticate users based on their MAC address. With this feature, the access point keeps an internal table of approved MAC addresses. When a wireless user tries to associate to the access point, the MAC address of the client must be on the approved list, or the association will be denied. Alternately, the AP may keep a table of known "bad" MAC addresses, and permit all devices that are not on the list.

Unfortunately, this is not an ideal security mechanism. Maintaining MAC tables on every device can be cumbersome, requiring all client devices to have their MAC addresses recorded and uploaded to the APs. Even worse, MAC addresses can often be changed in software. By observing MAC addresses in use on a wireless network, a determined attacker can **spoof** (impersonate) an approved MAC address and successfully associate to the AP. While MAC filtering will prevent unintentional users and even most curious individuals from accessing the network, MAC filtering alone cannot prevent attacks from determined attackers. MAC filters are useful for temporarily limiting access from misbehaving clients. For example, if a laptop has a virus that sends large amounts of spam or other traffic, its MAC address can be added to the filter table to stop the traffic immediately. This will buy you time to track down the user and fix the problem.

Another popular authentication feature of wireless the so-called *closed net-work*. In a typical network, APs will broadcast their ESSID many times per second, allowing wireless clients (as well as tools such as NetStumbler) to find the network and display its presence to the user. In a closed network, the AP does not beacon the ESSID, and users must know the full name of the network before the AP will allow association. This prevents casual users from discovering the network and selecting it in their wireless client.

There are a number of drawbacks to this feature. Forcing users to type in the full ESSID before connecting to the network is error prone and often leads to support calls and complaints. Since the network isn't obviously present in site survey tools like NetStumbler, this can prevent your networks from showing up on war driving maps. But it also means that other network builders cannot easily find your network either, and specifically won't know that you are already using a given channel. A conscientious neighbor may perform a site survey, see no nearby networks, and install their own network on the same channel you are using. This will cause interference problems for both you and your neighbor.

Finally, using closed networks ultimately adds little to your overall networks security. By using passive monitoring tools (such as Kismet), a skilled user can detect frames sent from your legitimate clients to the AP. These frames necessarily contain the network name. A malicious user can then use this name to associate to the access point, just like a normal user would.

Encryption is probably the best tool we have for authenticating wireless users. Through strong encryption, we can uniquely identify a user in a manner that is very difficult to spoof, and use that identity to determine further network access. Encryption also has the benefit of adding a layer of privacy by preventing eavesdroppers from easily watching network traffic.

The most widely employed encryption method on wireless networks is **WEP** encryption. WEP stands for wired equivalent privacy, and is supported by virtually all 802.11a/b/g equipment. WEP uses a shared 40-bit key to encrypt data between the access point and client. The key must be entered on the APs as well as on each of the clients. With WEP enabled, wireless clients cannot associate with the AP until they use the correct key. An eavesdropper listening to a WEP-enabled network will still see traffic and MAC addresses, but the data payload of each packet is encrypted. This provides a fairly good authentication mechanism while also adding a bit of privacy to the network.

WEP is definitely not the strongest encryption solution available. For one thing, the WEP key is shared between all users. If the key is compromised (say, if one user tells a friend what the password is, or if an employee is let go) then changing the password can be prohibitively difficult, since all APs and client devices need to be changed. This also means that legitimate users of the network can still eavesdrop on each others' traffic, since they all know the shared key.

The key itself is often poorly chosen, making offline cracking attempts feasible. Even worse, the implementation of WEP itself is broken in many access points, making it even easier to crack some networks. While manufacturers have implemented a number of extensions to WEP (such as longer keys and fast rotation schemes), these extensions are not part of the standard, and generally will not interoperate between equipment from different manufacturers. By upgrading to the most recent firmware for all of your wireless devices, you can prevent some of the early attacks found in WEP.

WEP can still be a useful authentication tool. Assuming your users can be trusted not to give away the password, you can be fairly sure that your wireless clients are legitimate. While WEP cracking is possible, it is beyond the skill of most users. WEP is quite useful for securing long distance point-topoint links, even on generally open networks. By using WEP on such a link, you will discourage others from associating to the link, and they will likely use other available APs instead. Think of WEP as a handy "keep out" sign for your network. Anyone who detects the network will see that a key is required, making it clear that they are not welcome to use it.

WEPs greatest strength is its interoperability. In order to comply with the 802.11 standards, all wireless devices support basic WEP. While it isn't the strongest method available, it is certainly the most commonly implemented encryption feature. We will look at other more advanced encryption techniques later in this chapter.

For more details about the state of WEP encryption, see these papers:

- · http://www.isaac.cs.berkeley.edu/isaac/wep-faq.html
- http://www.cs.umd.edu/~waa/wireless.pdf
- http://www.crypto.com/papers/others/rc4\_ksaproc.ps

Another data-link layer authentication protocol is *Wi-Fi Protected Access*, or *WPA*. WPA was created specifically to deal with the known problems with WEP mentioned earlier. It provides a significantly stronger encryption scheme, and can use a shared private key, unique keys assigned to each user, or even SSL certificates to authenticate both the client and the access point. Authentication credentials are checked using the 802.1X protocol,

which can consult a third party database such as RADIUS. Through the use of *Temporal Key Integrity Protocol* (*TKIP*), keys can be rotated quickly over time, further reducing the likelihood that a particular session can be cracked. Overall, WPA provides significantly better authentication and privacy than standard WEP.

WPA requires fairly recent access point hardware and up-to-date firmware on all wireless clients, as well as a substantial amount of configuration. If you are installing a network in a setting where you control the entire hardware platform, WPA can be ideal. By authenticating both clients and APs, it solves the rogue access point problem and provides many significant advantages over WEP. But in most network settings where the vintage of hardware is mixed and the knowledge of wireless users is limited, WPA can be a nightmare to install. It is for this reason that most sites continue to use WEP, if encryption is used at all.

# Captive portals

One common authentication tool used on wireless networks is the *captive portal*. A captive portal uses a standard web browser to give a wireless user the opportunity to present login credentials. It can also be used to present information (such as an Acceptable Use Policy) to the user before granting further access. By using a web browser instead of a custom program for authentication, captive portals work with virtually all laptops and operating systems. Captive portals are typically used on open networks with no other authentication methods (such as WEP or MAC filters).

To begin, a wireless user opens their laptop and selects the network. Their computer requests a DHCP lease, which is granted. They then use their web browser to go to any site on the Internet.

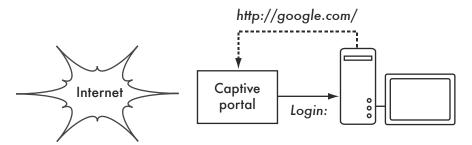


Figure 6.1: The user requests a web page and is redirected.

Instead of receiving the requested page, the user is presented with a login screen. This page can require the user to enter a user name and password, simply click a "login" button, type in numbers from a pre-paid ticket, or enter any other credentials that the network administrators require. The user then

enters their credentials, which are checked by the access point or another server on the network. All other network access is blocked until these credentials are verified.

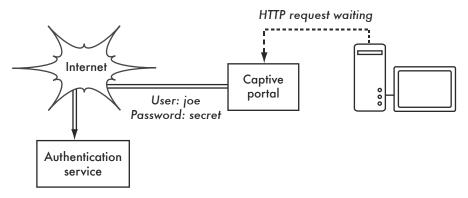


Figure 6.2: The user's credentials are verified before further network access is granted. The authentication server can be the access point itself, another machine on the local network, or a server anywhere on the Internet.

Once authenticated, the user is permitted to access network resources, and is typically redirected to the site they originally requested.

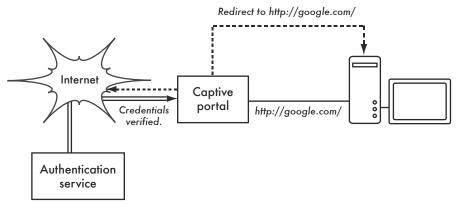


Figure 6.3: After authenticating, the user is permitted to access the rest of the network.

Captive portals provide no encryption for the wireless users, instead relying on the MAC and IP address of the client as a unique identifier. Since this is not necessarily very secure, many implementations will require the user to re-authenticate periodically. This can often be automatically done by minimizing a special pop-up browser window when the user first logs in.

Since they do not provide strong encryption, captive portals are not a very good choice for networks that need to be locked down to only allow access

from trusted users. They are much more suited to cafes, hotels, and other public access locations where casual network users are expected.

In public or semi-public network settings, encryption techniques such as WEP and WPA are effectively useless. There is simply no way to distribute public or shared keys to members of the general public without compromising the security of those keys. In these settings, a simple application such as a captive portal provides a level of service somewhere between completely open and completely closed.

#### Popular hotspot projects

- Chillispot (http://www.chillispot.info/). Chillispot is a captive portal designed to authenticate against an existing user credentials database, such as RADUIS. Combined with the application phpMyPrePaid, pre-paid ticket based authentication can be implemented very easily You can download phpMyPrePaid from http://sourceforge.net/projects/phpmyprepaid/.
- WiFi Dog (*http://www.wifidog.org/*). WiFi Dog provides a very complete captive portal authentication package in very little space (typically under 30kb). From a user's perspective, it requires no pop-up or javascript support, allowing it to work on a wider variety of wireless devices.
- m0n0wall (http://m0n0.ch/wall/). m0n0wall is a complete embedded operating system based on FreeBSD. It includes a captive portal with RADIUS support, as well as a PHP web server.
- NoCatSplash (http://nocat.net/download/NoCatSplash/) provides a customizable splash page to your users, requiring them to click a "login" button before using the network. This is useful for identifying the operators of the network and displaying rules for network access. It provides a very easy solution in situations where you need to provide users of an open network with information and an acceptable use policy.

# Privacy

Most users are blissfully unaware that their private email, chat conversations, and even passwords are often sent "in the clear" over dozens of untrusted networks before arriving at their ultimate destination on the Internet. However mistaken they may be, users still typically have some expectation of privacy when using computer networks.

Privacy can be achieved, even on untrusted networks such as public access points and the Internet. The only proven effective method for protecting privacy is the use of strong *end-to-end encryption*.

Encryption techniques such as WEP and WPA attempt to address the privacy issue at layer two, the data-link layer. This does protect against eavesdroppers listening in on the wireless connection, but this protection ends at the access point. If the wireless client uses insecure protocols (such as POP or simple SMTP for receiving and sending email), then users beyond the AP can still log the session and see the sensitive data. As mentioned earlier, WEP also suffers from the fact that it uses a shared private key. This means that legitimate wireless users can eavesdrop on each other, since they all know the private key.

By using encryption to the remote end of the connection, users can neatly sidestep the entire problem. These techniques work well even on untrusted public networks, where eavesdroppers are listening and possibly even manipulating data coming from the access point.

To ensure data privacy, good end-to-end encryption should provide the following features:

- Verified authentication of the remote end. The user should be able to know without a doubt that the remote end is who it claims to be. Without authentication, a user could give sensitive data to anyone claiming to be the legitimate service.
- **Strong encryption methods**. The encryption algorithm should stand up to public scrutiny, and not be easily decrypted by a third party. There is no security in obscurity, and strong encryption is even stronger when the algorithm is widely known and subject to peer review. A good algorithm with a suitably large and protected key can provide encryption that is unlikely to be broken by any effort in our lifetimes using current technology.
- **Public key cryptography**. While not an absolute requirement for end-toend encryption, the use of public key cryptography instead of a shared key can ensure that an individual's data remains private, even if the key of another user of the service is compromised. It also solves certain problems with distributing keys to users over untrusted networks.
- Data encapsulation. A good end-to-end encryption mechanism protects as much data as possible. This can range from encrypting a single email transaction to encapsulation of all IP traffic, including DNS lookups and other supporting protocols. Some encryption tools simply provide a secure channel that other applications can use. This allows users to run any program they like and still have the protection of strong encryption, even if the programs themselves don't support it.

Be aware that laws regarding the use of encryption vary widely from place to place. Some countries treat encryption as munitions, and may require a permit, escrow of private keys, or even prohibit its use altogether. Before

implementing any solution that involves encryption, be sure to verify that use of this technology is permitted in your local area.

In the following sections, we'll take a look at some specific tools that can provide good protection for your users' data.

# SSL

The most widely available end-to-end encryption technology is **Secure Sockets Layer**, known simply as **SSL**. Built into virtually all web browsers, SSL uses public key cryptography and a trusted **public key infrastructure** (**PKI**) to secure data communications on the web. Whenever you visit a web URL that starts with https, you are using SSL.

The SSL implementation built into web browsers includes a collection of certificates from trusted sources, called *certificate authorities* (*CA*). These certificates are cryptographic keys that are used to verify the authenticity of websites. When you browse to a website that uses SSL, the browser and the server first exchange certificates. The browser then verifies that the certificate provided by the server matches its DNS host name, that it has not expired, and that it is signed by a trusted certificate authority. The server optionally verifies the identity of the browser's certificate. If the certificates are approved, the browser and server then negotiate a master session key using the previously exchanged certificates to protect it. That key is then used to encrypt all communications until the browser disconnects. This kind of data encapsulation is known as a *tunnel*.

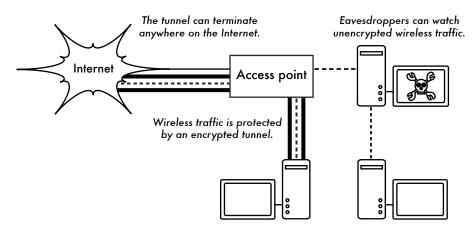


Figure 6.4: Eavesdroppers must break strong encryption to monitor traffic over an encrypted tunnel. The conversation inside the tunnel is identical to any other unencrypted conversation.

The use of certificates with a PKI not only protects the communication from eavesdroppers, but also prevents so-called *man-in-the-middle* (*MITM*) at-

tacks. In a man-in-the-middle attack, a malicious user intercepts all communication between the browser and the server. By presenting counterfeit certificates to both the browser and the server, the malicious user could carry on two simultaneous encrypted sessions. Since the malicious user knows the secret on both connections, it is trivial to observe and manipulate data passing between the server and the browser.

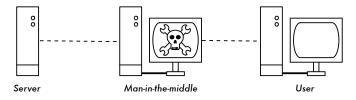


Figure 6.5: The man-in-the-middle effectively controls everything the user sees, and can record and manipulate all traffic. Without a public key infrastructure to verify the authenticity of keys, strong encryption alone cannot protect against this kind of attack.

Use of a good PKI prevents this kind of attack. In order to be successful, the malicious user would have to present a certificate to the client that is signed by a trusted certificate authority. Unless a CA has been compromised (very unlikely) or the user is tricked into accepting the forged certificate, then such an attack is not possible. This is why it is vitally important that users understand that ignoring warnings about expired or improper certificates is very dangerous, especially when using wireless networks. By clicking the "ignore" button when prompted by their browser, users open themselves up to many potential attacks.

SSL is not only used for web browsing. Insecure email protocols such as IMAP, POP, and SMTP can be secured by wrapping them in an SSL tunnel. Most modern email clients support IMAPS and POPS (secure IMAP and POP) as well as SSL/TLS protected SMTP. If your email server does not provide SSL support, you can still secure it with SSL using a package like Stunnel (*http://www.stunnel.org/*). SSL can be used to effectively secure just about any service that runs over TCP.

# SSH

Most people think of SSH as a secure replacement for telnet, just as scp and sftp are the secure counterparts of rcp and ftp. But SSH is much more than encrypted remote shell. Like SSL, it uses strong public key cryptography to verify the remote server and encrypt data. Instead of a PKI, it uses a key fingerprint cache that is checked before a connection is permitted. It can use passwords, public keys, or other methods for user authentication.

Many people do not know that SSH can also act as a general purpose encrypting tunnel, or even an encrypting web proxy. By first establishing an SSH connection to a trusted location near (or even on) a remote server, insecure protocols can be protected from eavesdropping and attack.

While this technique may be a bit advanced for many users, network architects can use SSH to encrypt traffic across untrusted links, such as wireless point-to-point links. Since the tools are freely available and run over standard TCP, any educated user can implement SSH connections for themselves, providing their own end-to-end encryption without administrator intervention.

OpenSSH (*http://openssh.org/*) is probably the most popular implementation on Unix-like platforms. Free implementations such as Putty (*http://www.putty.nl/*) and WinSCP (*http://winscp.net/*) are available for Windows. OpenSSH will also run on Windows under the Cygwin package (*http://www.cygwin.com/*). These examples will assume that you are using a recent version of OpenSSH.

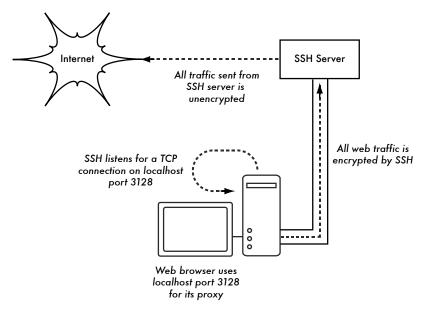


Figure 6.6: The SSH tunnel protects web traffic up to the SSH server itself.

To establish an encrypted tunnel from a port on the local machine to a port on the remote side, use the **-L** switch. For example, suppose you want to forward web proxy traffic over an encrypted link to the squid server at *squid.example.net*. Forward port 3128 (the default proxy port) using this command:

ssh -fN -g -L3128:squid.example.net:3128 squid.example.net

The **-fN** switches instruct ssh to fork into the background after connecting. The **-g** switch allows other users on your local segment to connect to the local machine and use it for encryption over the untrusted link. OpenSSH will use a public key for authentication if you have set one up, or it will prompt you for your password on the remote side. You can then configure your web browser to connect to localhost port 3128 as its web proxy service. All web traffic will then be encrypted before transmission to the remote side.

SSH can also act as a dynamic SOCKS4 or SOCKS5 proxy. This allows you to create an encrypting web proxy, without the need to set up squid. Note that this is not a caching proxy; it simply encrypts all traffic.

ssh -fN -D 8080 remote.example.net

Configure your web browser to use SOCKS4 or SOCKS5 on local port 8080, and away you go.

SSH can encrypt data on any TCP port, including ports used for email. It can even compress the data along the way, which can decrease latency on low capacity links.

ssh -fNCg -L110:localhost:110 -L25:localhost:25 mailhost.example.net

The **-C** switch turns on compression. You can add as many port forwarding rules as you like by specifying the **-L** switch multiple times. Note that in order to bind to a local port less than 1024, you must have root privileges on the local machine.

These are just a few examples of the flexibility of SSH. By implementing public keys and using the ssh forwarding agent, you can automate the creation of encrypted tunnels throughout your wireless network, and protect your communications with strong encryption and authentication.

# OpenVPN

OpenVPN is a free, open source VPN implementation built on SSL encryption. There are OpenVPN client implementations for a wide range of operating systems, including Linux, Windows 2000/XP and higher, OpenBSD, FreeBSD, NetBSD, Mac OS X, and Solaris. Being a VPN, it encapsulates all traffic (including DNS and all other protocols) in an encrypted tunnel, not just a single TCP port. Most people find it considerably easier to understand and configure than IPSEC.

OpenVPN also has some disadvantages, such as fairly high latency. Some amount of latency is unavoidable since all encryption/decryption is done in user space, but using relatively new computers on either end of the tunnel can minimize this. While it can use traditional shared keys, OpenVPN really shines when used with SSL certificates and a certificate authority. OpenVPN has many advantages that make it a good option for providing end-to-end security.

Some of these reasons include:

- It is based on a proven, robust encryption protocol (SSL and RSA)
- · It is relatively easy to configure
- · It functions across many different platforms
- · It is well documented
- It's free and open source.

OpenVPN needs to connect to a single TCP or UDP port on the remote side. Once established, it can encapsulate all data down to the Networking layer, or even down to the Data-Link layer, if your solution requires it. You can use it to create robust VPN connections between individual machines, or simply use it to connect network routers over untrusted wireless networks.

VPN technology is a complex field, and is a bit beyond the scope of this section to go into more detail. It is important to understand how VPNs fit into the structure of your network in order to provide the best possible protection without opening up your organization to unintentional problems. There are many good online resources that deal with installing OpenVPN on a server and client, we recommend this article from Linux Journal: *http://www.linuxjournal.com/article/7949* as well as the official HOWTO: *http://openvpn.net/howto.html* 

# Tor & Anonymizers

The Internet is basically an open network based on trust. When you connect to a web server across the Internet, your traffic passes through many different routers, owned by a great variety of institutions, corporations and individuals. In principle, any one of these routers has the ability to look closely at your data, seeing the source and destination addresses, and quite often also the actual content of the data. Even if your data is encrypted using a secure protocol, it is possible for your Internet provider to monitor the amount of data transferred, as well as the source and destination of that data. Often this is enough to piece together a fairly complete picture of your activities on-line.

Privacy and anonymity are important, and closely linked to each other. There are many valid reasons to consider protecting your privacy by **anonymizing** your network traffic. Suppose you want to offer Internet connectivity to your local community by setting up a number of access points for people to connect to. Whether you charge them for their access or not, there is always the

risk that people use the network for something that is not legal in your country or region. You could plead with the legal system that this particular illegal action was not performed by yourself, but could have been performed by anyone connecting to your network. The problem is neatly sidestepped if it were technically infeasible to determine where your traffic was actually headed. And what about on-line censorship? Publishing web pages anonymously may also be necessary to avoid government censorship.

There are tools that allow you to anonymize your traffic in relatively easy ways. The combination of **Tor** (*http://www.torproject.org/*) and **Privoxy** (*http://www.privoxy.org/*) is a powerful way to run a local proxy server that will pass your Internet traffic through a number of servers all across the net, making it very difficult to follow the trail of information. Tor can be run on a local PC, under Microsoft Windows, Mac OSX, Linux and a variety of BSD's, where it anonymizes traffic from the browser on that particular machine. Tor and Privoxy can also be installed on a gateway server, or even a small embedded access point (such as a Linksys WRT54G) where they provides anonymity to all network users automatically.

Tor works by repeatedly bouncing your TCP connections across a number of servers spread throughout the Internet, and by wrapping routing information in a number of encrypted layers (hence the term **onion routing**), that get peeled off as the packet moves across the network. This means that, at any given point in the network, the source and destination addresses cannot be linked together. This makes traffic analysis extremely difficult.

The need for the Privoxy privacy proxy in connection with Tor is due to the fact that name server queries (DNS queries) in most cases are not passed through the proxy server, and someone analyzing your traffic would easily be able to see that you were trying to reach a specific site (say *google.com*) by the fact that you sent a DNS query to translate google.com to the appropriate IP address. Privoxy connects to Tor as a SOCKS4a proxy, which uses host-names (not IP addresses) to get your packets to the intended destination.

In other words, using Privoxy with Tor is a simple and effective way to prevent traffic analysis from linking your IP address with the services you use online. Combined with secure, encrypted protocols (such as those we have seen in this chapter), Tor and Privoxy provide a high level of anonymity on the Internet.

# Network Monitoring

Network monitoring is the use of logging and analysis tools to accurately determine traffic flows, utilization, and other performance indicators on a network. Good monitoring tools give you both hard numbers and graphical aggregate representations of the state of the network. This helps you to visualize precisely what is happening, so you know where adjustments may be needed. These tools can help you answer critical questions, such as:

- · What are the most popular services used on the network?
- · Who are the heaviest network users?
- · What other wireless channels are in use in my area?
- · Are users installing wireless access points on my private wired network?
- · At what time of the day is the network most utilized?
- · What sites do your users frequent?
- Is the amount of inbound or outbound traffic close to our available network capacity?
- Are there indications of an unusual network situation that is consuming bandwidth or causing other problems?
- Is our Internet Service Provider (ISP) providing the level of service that we are paying for? This should be answered in terms of available bandwidth, packet loss, latency, and overall availability.

And perhaps the most important question of all:

· Do the observed traffic patterns fit our expectations?

Let's look at how a typical system administrator can make good use of network monitoring tools.

# An effective network monitoring example

For the purposes of example, let's assume that we are in charge of a network that has been running for three months. It consists of 50 computers and three servers: email, web, and proxy servers. While initially things are going well, users begin to complain of slow network speeds and an increase in spam emails. As time goes on, computer performance slows to a crawl (even when not using the network), causing considerable frustration in your users.

With frequent complaints and very low computer usage, the Board is questioning the need for so much network hardware. The Board also wants evidence that the bandwidth they are paying for is actually being used. As the network administrator, you are on the receiving end of these complaints. How can you diagnose the sudden drop in network and computer performance and also justify the network hardware and bandwidth costs?

## Monitoring the LAN (local traffic)

To get an idea of exactly what is causing the slow down, you should begin by looking at traffic on the local LAN. There are several advantages to monitoring local traffic:

- Troubleshooting is greatly simplified.
- Viruses can be detected and eliminated.
- · Malicious users can be detected and dealt with.
- Network hardware and resources can be justified with real statistics.

Assume that all of the switches support the *Simple Network Management Protocol* (*SNMP*). SNMP is an application-layer protocol designed to facilitate the exchange of management information between network devices. By assigning an IP address to each switch, you are able to monitor all the interfaces on that switch, observing the entire network from a single point. This is much easier than enabling SNMP on all computers in a network.

By using a free tool such as MRTG (see **Page 190**), you can monitor each port on the switch and present data graphically, as an aggregate average over time. The graphs are accessible from the web, so you are able to view the graphs from any machine at anytime.

With MRTG monitoring in place, it becomes obvious that the internal LAN is swamped with far more traffic than the Internet connection can support, even when the lab is unoccupied. This is a pretty clear indication that some of the computers are infested with a network virus. After installing good anti-virus and anti-spyware software on all of the machines, the internal LAN traffic settles down to expected levels. The machines run much more quickly, spam emails are reduced, and the users' morale quickly improves.

#### Monitoring the WAN (external traffic)

In addition to watching the traffic on the internal LAN, you need to demonstrate that the bandwidth the organization is paying for is actually what they are getting from their ISP. You can achieve this by monitoring *external traffic*.

External traffic is generally classified as anything sent over a *Wide Area Network* (*WAN*). Anything received from (or sent to) a network other than your internal LAN also qualifies as external traffic. The advantages of monitoring external traffic include:

• Internet bandwidth costs are justified by showing actual usage, and whether that usage agrees with your ISP's bandwidth charges.

- Future capacity needs are estimated by watching usage trends and predicting likely growth patterns.
- Intruders from the Internet are detected and filtered before they can cause problems.

Monitoring this traffic is easily done with the use of MRTG on an SNMP enabled device, such as a router. If your router does not support SNMP, then you can add a switch between your router and your ISP connection, and monitor the port traffic just as you would with an internal LAN.

#### **Detecting Network Outages**

With monitoring tools in place, you now have an accurate measurement of how much bandwidth the organization is using. This measurement should agree with your ISP's bandwidth charges. It can also indicate the actual throughput of your connection if you are using close to your available capacity at peak times. A "flat top" graph is a fairly clear indication that you are operating at full capacity. **Figure 6.7** shows flat tops in peak outbound traffic in the middle of every day except Sunday.

It is clear that your current Internet connection is overutilized at peak times, causing network lag. After presenting this information to the Board, you can make a plan for further optimizing your existing connection (by upgrading your proxy server and using other techniques in this book) and estimate how soon you will need to upgrade your connection to keep up with the demand. This is also an excellent time to review your operational policy with the Board, and discuss ways to bring actual usage in line with that policy.

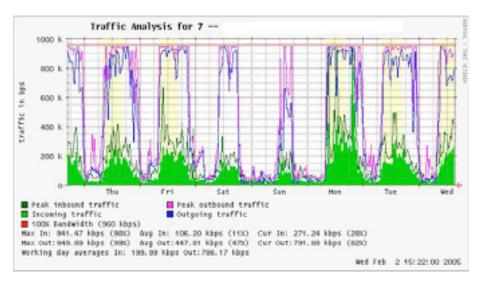


Figure 6.7: A graph with a "flat top" is one indication of overutilization.

Later in the week, you receive an emergency phone call in the evening. Apparently, no one in the lab can browse the web or send email. You rush to the lab and hastily reboot the proxy server, with no results. Browsing and email are still broken. You then reboot the router, but there is still no success. You continue eliminating the possible fault areas one by one until you realize that the network switch is off - a loose power cable is to blame. After applying power, the network comes to life again.

How can you troubleshoot such an outage without such time consuming trial and error? Is it possible to be notified of outages as they occur, rather than waiting for a user to complain? One way to do this is to use a program such as **Nagios** that continually polls network devices and notifies you of outages. Nagios will report on the availability of various machines and services, and will alert you to machines that have gone down. In addition to displaying the network status graphically on a web page, it will send notifications via SMS or email, alerting you immediately when problems arise.

With good monitoring tools in place, you will be able to justify the cost of equipment and bandwidth by effectively demonstrating how it is being used by the organization. You are notified automatically when problems arise, and you have historical statistics of how the network devices are performing. You can check the current performance against this history to find unusual behavior, and head off problems before they become critical. When problems do come up, it is simple to determine the source and nature of the problem. Your job is easier, the Board is satisfied, and your users are much happier.

# Monitoring your network

Managing a network without monitoring is similar to driving a vehicle without a speedometer or a fuel gauge, with your eyes closed. How do you know how fast you are going? Is the car consuming fuel as efficiently as promised by the dealers? If you do an engine overhaul several months later, is the car any faster or more efficient than it was before?

Similarly, how can you pay for an electricity or water bill without seeing your monthly usage from a meter? You must have an account of your network bandwidth utilization in order to justify the cost of services and hardware purchases, and to account for usage trends.

There are several benefits to implementing a good monitoring system for your network:

1. **Network budget and resources are justified.** Good monitoring tools can demonstrate without a doubt that the network infrastructure (bandwidth, hardware, and software) is suitable and able to handle the requirements of network users.

- Network intruders are detected and filtered. By watching your network traffic, you can detect attackers and prevent access to critical internal servers and services.
- 3. **Network viruses are easily detected.** You can be alerted to the presence of network viruses, and take appropriate action before they consume Internet bandwidth and destabilize your network
- 4. **Troubleshooting of network problems is greatly simplified.** Rather than attempting "trial and error" to debug network problems, you can be instantly notified of specific problems. Some kinds of problems can even be repaired automatically.
- 5. **Network performance can be highly optimized.** Without effective monitoring, it is impossible to fine tune your devices and protocols to achieve the best possible performance.
- 6. **Capacity planning is much easier.** With solid historical performance records, you do not have to "guess" how much bandwidth you will need as your network grows.
- 7. **Proper network usage can be enforced.** When bandwidth is a scarce resource, the only way to be fair to all users is to ensure that the network is being used for its intended purpose.

Fortunately, network monitoring does not need to be an expensive undertaking. There are many freely available open source tools that will show you exactly what is happening on your network in considerable detail. This section will help you identify many invaluable tools and how best to use them.

# The dedicated monitoring server

While monitoring services can be added to an existing network server, it is often desirable to dedicate one machine (or more, if necessary) to network monitoring. Some applications (such as *ntop*) require considerable resources to run, particularly on a busy network. But most logging and monitoring programs have modest RAM and storage requirements, typically with little CPU power required. Since open source operating systems (such as Linux or BSD) make very efficient use of hardware resources, this makes it possible to build a very capable monitoring server from recycled PC parts. There is usually no need to purchase a brand new server to relegate to monitoring duties.

The exception to this rule is in very large installations. If your network includes more than a few hundred nodes, or if you consume more than 50 Mbps of Internet bandwidth, you will likely need to split up monitoring duties between a few dedicated machines. This depends largely on exactly what you want to monitor. If you are attempting to account for all services accessed per MAC address, this will consume considerably more resources than simply measuring network flows on a switch port. But for the majority of installations, a single dedicated monitoring machine is usually enough.

While consolidating monitoring services to a single machine will streamline administration and upgrades, it can also ensure better ongoing monitoring. For example, if you install monitoring services on a web server, and that web server develops problems, then your network may not be monitored until the problem is resolved.

To a network administrator, the data collected about network performance is nearly as important as the network itself. Your monitoring should be robust and protected from service outages as well as possible. Without network statistics, you are effectively blind to problems with the network.

#### Where does the server fit in my network?

If you are only interested in collecting network flow statistics from a router, you can do this from just about anywhere on the LAN. This provides simple feedback about utilization, but cannot give you comprehensive details about usage patterns. **Figure 6.8** shows a typical MRTG graph generated from the Internet router. While the inbound and outbound utilization are clear, there is no detail about which computers, users, or protocols are using bandwidth.

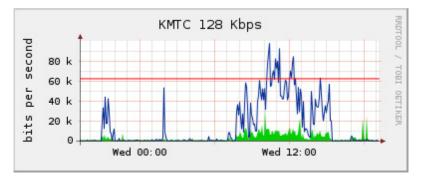


Figure 6.8: Polling the edge router can show you the overall network utilization, but you cannot break the data down further into machines, services, and users.

For more detail, the dedicated monitoring server must have access to everything that needs to be watched. Typically, this means it must have access to the entire network. To monitor a WAN connection, such as the Internet link to your ISP, the monitoring server must be able to see the traffic passing through the edge router. To monitor a LAN, the monitoring server is typically connected to a *monitor port* on the switch. If multiple switches are used in an installation, the monitoring server may need a connection to all of them. That connection can either be a physical cable, or if

your network switches support it, a VLAN specifically configured for monitoring traffic.

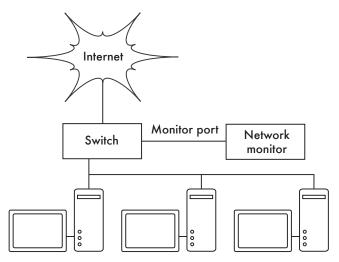


Figure 6.9: Use the monitor port on your switch to observe traffic crossing all of the network ports.

If monitor port functionality is not available on your switch, the monitoring server may be installed between your internal LAN and the Internet. While this will work, it introduces a single point of failure for the network, as the network will fail if the monitoring server develops a problem. It is also a potential performance bottleneck, if the server cannot keep up with the demands of the network.

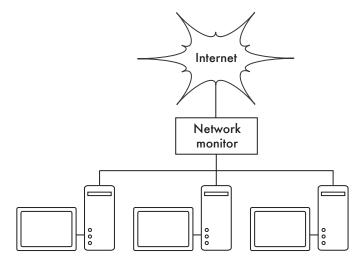


Figure 6.10: By inserting a network monitor between the LAN and your Internet connection, you can observe all network traffic.

A better solution is to use a simple network hub (not a switch) which connects the monitoring machine to the internal LAN, external router, and the monitoring machine. While this does still introduce an additional point of failure to the network (since the entire network will be unreachable if the hub dies), hubs are generally considered to be much more reliable than routers. They are also very easily replaced should they fail.

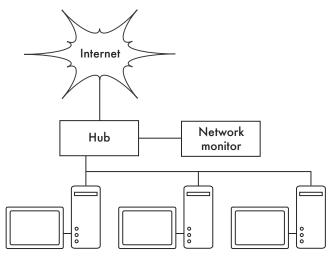


Figure 6.11: If your switch does not provide monitor port functionality, you can insert a network hub between your Internet router and the LAN, and connect the monitoring server to the hub.

Once your monitoring server is in place, you are ready to start collecting data.

## What to monitor

It is possible to plot just about any network event and watch its value on a graph over time. Since every network is slightly different, you will have to decide what information is important in order to gauge the performance of your network.

Here are some important indicators that many network administrators will typically track.

#### Wireless statistics

- · Received signal and noise from all backbone nodes
- · Number of associated stations
- · Detected adjacent networks and channels
- Excessive retransmissions

· Radio data rate, if using automatic rate scaling

#### Switch statistics

- · Bandwidth usage per switch port
- · Bandwidth usage broken down by protocol
- · Bandwidth usage broken down by MAC address
- · Broadcasts as a percentage of total packets
- Packet loss and error rate

#### Internet statistics

- · Internet bandwidth use by host and protocol
- · Proxy server cache hits
- Top 100 sites accessed
- DNS requests
- · Number of inbound emails / spam emails / email bounces
- · Outbound email queue size
- · Availability of critical services (web servers, email servers, etc.).
- · Ping times and packet loss rates to your ISP
- Status of backups

#### System health statistics

- Memory usage
- Swap file usage
- Process count / zombie processes
- System load
- · Uninterruptible Power Supply (UPS) voltage and load
- · Temperature, fan speed, and system voltages
- Disk SMART status
- RAID array status

You should use this list as a suggestion of where to begin. As your network matures, you will likely find new key indicators of network performance, and you should of course track those as well. There are many freely available

tools that will show you as much detail as you like about what is happening on your network. You should consider monitoring the availability of any resource where unavailability would adversely affect your network users.

For example, your users may dial into modems on your site to gain remote access to your network. If all the modems are used, or if any are faulty, then users will be denied access and will probably complain. You can predict and avoid such problems by monitoring the number of available modems, and provisioning extra capacity before you run out.

Don't forget to monitor the monitoring machine itself, for example its CPU usage and disk space, in order to receive advance warning if it becomes overloaded or faulty. A monitoring machine that is low on resources can affect your ability to monitor the network effectively.

# Types of monitoring tools

We will now look at several different classes of monitoring tools. Network detection tools listen for the beacons sent by wireless access points, and display information such as the network name, received signal strength, and channel. Spot check tools are designed for troubleshooting and normally run interactively for short periods of time. A program such as ping may be considered an active spot check tool, since it generates traffic by polling a particular machine. Passive spot check tools include protocol analyzers, which inspect every packet on the network and provide complete detail about any network conversation (including source and destination addresses, protocol information, and even application data). Trending tools perform unattended monitoring over long periods, and typically plot the results on a graph. Realtime monitoring tools perform similar monitoring, but notify administrators immediately if they detect a problem. Throughput testing tools tell you the actual bandwidth available between two points on a network. Intrusion detection tools watch for undesirable or unexpected network traffic, and take appropriate action (typically denying access and/or notifying a network administrator). Finally, benchmarking tools estimate the maximum performance of a service or network connection.

# Network detection

The simplest wireless monitoring tools simply provide a list of available networks, along with basic information (such as signal strength and channel). They let you quickly detect nearby networks and determine if they are in range or are causing interference.

• **The built-in client.** All modern operating systems provide built-in support for wireless networking. This typically includes the ability to scan for available networks, allowing the user to choose a network from a list. While

virtually all wireless devices are guaranteed to have a simple scanning utility, functionality can vary widely between implementations. These tools are typically only useful for configuring a computer in a home or office setting. They tend to provide little information apart from network names and the available signal to the access point currently in use.

- Netstumbler (http://www.netstumbler.com/). This is the most popular tool for detecting wireless networks using Microsoft Windows. It supports a variety of wireless cards, and is very easy to use. It will detect open and encrypted networks, but cannot detect "closed" wireless networks. It also features a signal/noise meter that plots radio receiver data as a graph over time. It also integrates with a variety of GPS devices, for logging precise location and signal strength information. This makes Netstumbler a handy tool to have for an informal site survey.
- **Ministumbler** (*http://www.netstumbler.com/*). From the makers of Netstumbler, Ministumbler provides much of the same functionality as the Windows version, but works on the Pocket PC platform. Ministumbler is handy to run on a handheld PDA with a wireless card for detecting access points in the field.
- **Macstumbler** (*http://www.macstumbler.com/*). While not directly related to the Netstumbler, Macstumbler provides much of the same functionality but for the Mac OS X platform. It works with all Apple Airport cards.
- Wellenreiter (*http://www.wellenreiter.net/*). Wellenreiter is a nice graphical wireless network detector for Linux. It requires Perl and GTK, and supports Prism2, Lucent, and Cisco wireless cards.

# Spot check tools

What do you do when the network breaks? If you can't access a web page or email server, and clicking the reload button doesn't fix the problem, then you'll need to be able to isolate the exact location of the problem. These tools will help you to determine just where a connection problem exists.

This section is simply an introduction to commonly used troubleshooting tools. For more discussion of common network problems and how to diagnose them, see **Chapter 9, Troubleshooting**.

#### ping

Just about every operating system (including Windows, Mac OS X, and of course Linux and BSD) includes a version of the **ping** utility. It uses ICMP packets to attempt to contact a specified host, and tells you how long it takes to get a response.

Knowing what to ping is just as important as knowing how to ping. If you find that you cannot connect to a particular service in your web browser (say, *http://yahoo.com/*), you could try to ping it:

#### \$ ping yahoo.com

PING yahoo.com (66.94.234.13): 56 data bytes 64 bytes from 66.94.234.13: icmp\_seq=0 ttl=57 time=29.375 ms 64 bytes from 66.94.234.13: icmp\_seq=1 ttl=56 time=35.467 ms 64 bytes from 66.94.234.13: icmp\_seq=2 ttl=56 time=34.158 ms ^C --- yahoo.com ping statistics ---3 packets transmitted, 3 packets received, 0% packet loss round-trip min/avg/max/stddev = 29.375/33.000/35.467/2.618 ms

Hit control-C when you are finished collecting data. If packets take a long time to come back, there may be network congestion. If return ping packets have an unusually low *Time To Live (TTL)*, you may have routing problems between your machine and the remote end. But what if the ping doesn't return any data at all? If you are pinging a name instead of an IP address, you may be running into DNS problems.

Try pinging an IP address on the Internet. If you can't reach it, it's a good idea to see if you can ping your default router:

#### \$ ping 69.90.235.230

PING 69.90.235.230 (69.90.235.230): 56 data bytes 64 bytes from 69.90.235.230: icmp\_seq=0 ttl=126 time=12.991 ms 64 bytes from 69.90.235.230: icmp\_seq=1 ttl=126 time=14.869 ms 64 bytes from 69.90.235.230: icmp\_seq=2 ttl=126 time=13.897 ms ^C --- 216.231.38.1 ping statistics ---3 packets transmitted, 3 packets received, 0% packet loss round-trip min/avg/max/stddev = 12.991/13.919/14.869/0.767 ms

If you can't ping your default router, then chances are you won't be able to get to the Internet either. If you can't even ping other IP addresses on your local LAN, then it's time to check your connection. If you're using Ethernet, is it plugged in? If you're using wireless, are you connected to the proper wireless network, and is it in range?

Network debugging with ping is a bit of an art, but it is useful to learn. Since you will likely find ping on just about any machine you will work on, it's a good idea to learn how to use it well.

#### traceroute and mtr

*http://www.bitwizard.nl/mtr/*. As with ping, traceroute is found on most operating systems (it's called **tracert** in some versions of Microsoft Windows). By running traceroute, you can find the location of problems between your computer and any point on the Internet:

```
$ traceroute -n google.com
traceroute to google.com (72.14.207.99), 64 hops max, 40 byte packets
1 10.15.6.1 4.322 ms 1.763 ms 1.731 ms
2 216.231.38.1 36.187 ms 14.648 ms 13.561 ms
3 69.17.83.233 14.197 ms 13.256 ms 13.267 ms
4 69.17.83.150 32.478 ms 29.545 ms 27.494 ms
5 198.32.176.31 40.788 ms 28.160 ms 28.115 ms
6 66.249.94.14 28.601 ms 29.913 ms 28.811 ms
7 172.16.236.8 2328.809 ms 2528.944 ms 2428.719 ms
8 * * *
```

The **-n** switch tells traceroute not to bother resolving names in DNS, and makes the trace run more quickly. You can see that at hop seven, the round trip time shoots up to more than two seconds, while packets seem to be discarded at hop eight. This might indicate a problem at that point in the network. If this part of the network is in your control, it might be worth starting your troubleshooting effort there.

**My TraceRoute** (**mtr**) is a handy program that combines ping and traceroute into a single tool. By running mtr, you can get an ongoing average of latency and packet loss to a single host, instead of the momentary snapshot that ping and traceroute provide.

	My trac	eroute	[v0.6	9]										
tesla.rob.swn (0.0.0	.0) (tos=	0x0 psi	ze=64	bitpatS	un Jan	8 20	:01:26	2006						
Keys: Help Displa	y mode Restart	statis	tics	Order	of fie	lds	quit							
		Pack	ets	Pings										
Host		Loss%	Snt	Last	Avg	Best	Wrst	StDev						
<ol> <li>gremlin.rob.swn</li> </ol>		0.0%	4	1.9	2.0	1.7	2.6	0.4						
<ol><li>erl.seal.speakea</li></ol>	sy.net	0.0%	4	15.5	14.0	12.7	15.5	1.3						
3. 220.ge-0-1-0.cr2	.seal.speakeasy.	0.0%	4	11.0	11.7	10.7	14.0	1.6						
4. fe-0-3-0.cr2.sfc	1.speakeasy.net	0.0%	4	36.0	34.7	28.7	38.1	4.1						
5. bas1-m.pao.yahoo	.com	0.0%	4	27.9	29.6	27.9	33.0	2.4						
6. so-1-1-0.pat1.dc	e.yahoo.com	0.0%	4	89.7	91.0	89.7	93.0	1.4						
7. ae1.p400.msr1.dc	n.yahoo.com	0.0%	4	91.2	93.1	90.8	99.2	4.1						
8. ge5-2.bas1-m.dcn	.yahoo.com	0.0%	4	89.3	91.0	89.3	93.4	1.9						
9. w2.rc.vip.dcn.ya	hoo.com	0.0%	3	91.2	93.1	90.8	99.2	4.1						

The data will be continuously updated and averaged over time. As with ping, you should hit control-C when you are finished looking at the data. Note that you must have root privileges to run mtr.

While these tools will not revel precisely what is wrong with the network, they can give you enough information to know where to continue troubleshooting.

# Protocol analyzers

Network protocol analyzers provide a great deal of detail about information flowing through a network, by allowing you to inspect individual packets. For wired networks, you can inspect packets at the data-link layer or above. For wireless networks, you can inspect information all the way down to individual 802.11 frames. Here are several popular (and free) network protocol analyzers:

#### Kismet

*http://www.kismetwireless.net/.* Kismet is a powerful wireless protocol analyzer for many platforms including Linux, Mac OS X, and even the embedded OpenWRT Linux distribution. It works with any wireless card that supports passive monitor mode. In addition to basic network detection, Kismet will passively log all 802.11 frames to disk or to the network in standard PCAP format, for later analysis with tools like Ethereal. Kismet also features associated client information, AP hardware fingerprinting, Netstumbler detection, and GPS integration.

Since it is a passive network monitor, it can even detect "closed" wireless networks by analyzing traffic sent by wireless clients. You can run Kismet on several machines at once, and have them all report over the network back to a central user interface. This allows for wireless monitoring over a large area, such as a university or corporate campus.



Figure 6.12: Kismet running on a Nokia 770 Internet Tablet

Since Kismet uses the radio card's passive monitor mode, it does all of this without transmitting any data. Kismet is an invaluable tool for diagnosing wireless network problems.

## KisMAC

*http://kismac.macpirate.ch/*. Exclusively for the Mac OS X platform, KisMAC does much of what Kismet can do, but with a slick Mac OS X graphical interface. It is a passive scanner that will log data to disk in PCAP format compatible with Wireshark. It supports passive scanning with AirportExtreme cards as well as a variety of USB wireless adapters.

### tcpdump

*http://www.tcpdump.org/.* **tcpdump** is a command-line tool for monitoring network traffic. It does not have all the bells and whistles of wireshark but it does use fewer resources. Tcpdump can capture and display all network protocol information down to the link layer. It can show all of the packet headers and data received, or just the packets that match particular criteria. Packets captured with tcpdump can be loaded into wireshark for visual analysis and further diagnostics. This is very useful if you wish to monitor an interface on a remote system and bring the file back to your local machine for analysis. The tcpdump tool is available as a standard tool in Unix derivatives (Linux, BSD, and Mac OS X). There is also a Windows port called *WinDump* available at *http://www.winpcap.org/windump/.* 

# Wireshark

*http://www.wireshark.org/.* Formerly known as **Ethereal**, Wireshark is a free network protocol analyzer for Unix and Windows. It is billed as "The World's Most Popular Network Protocol Analyzer."

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Figure 6.13: Wireshark (formerly Ethereal) is a powerful network protocol analyzer that can show you as much detail as you like about any packet.

Wireshark allows you to examine data from a live network or from a capture file on disk, and interactively browse and sort the captured data. Both summary and detailed information is available for each packet, including the full header and data portions. Wireshark has several powerful features, including a rich display filter language and the ability to view the reconstructed stream of a TCP session.

It can be daunting to use for first time users or those that are not familiar with the OSI layers. It is typically used to isolate and analyze specific traffic to or from an IP address, but it can be also used as a general purpose fault finding tool. For example, a machine infected with a network worm or virus can be identified by looking for the machine that is send out the same sort of TCPIP packets to large groups of IP addresses.

# Trending tools

Trending tools are used to see how your network is used over a long period of time. They work by periodically monitoring your network activity, and displaying a summary in a human-readable form (such as a graph). Trending tools collect data as well as analyze and report on it.

Below are some examples of trending tools. Some of them need to be used in conjunction with each other, as they are not stand-alone programs.

### MRTG

*http://oss.oetiker.ch/mrtg/.* The *Multi Router Traffic Grapher* (*MRTG*) monitors the traffic load on network links using SNMP. MRTG generates graphs that provide a visual representation of inbound and outbound traffic. These are typically displayed on a web page.

MRTG can be a little confusing to set up, especially if you are not familiar with SNMP. But once it is installed, MRTG requires virtually no maintenance, unless you change something on the system that is being monitored (such as its IP address).

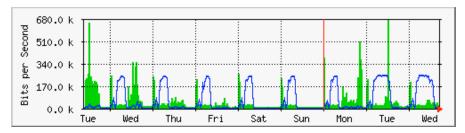


Figure 6.14: MRTG is probably the most widely installed network flow grapher.

## RRDtool

http://oss.oetiker.ch/rrdtool/. **RRD** is short for **Round Robin Database**. RRD is a database that stores information in a very compact way that does not

expand over time. *RRDtool* refers to a suite of tools that allow you to create and modify RRD databases, as well as generate useful graphs to present the data. It is used to keep track of time-series data (such as network bandwidth, machine room temperature, or server load average) and can display that data as an average over time.

Note that RRDtool itself does not contact network devices to retrieve data. It is merely a database manipulation tool. You can use a simple wrapper script (typically in shell or Perl) to do that work for you. RRDtool is also used by many full featured front-ends that present you with a friendly web interface for configuration and display. RRD graphs give you more control over display options and the number of items available on a graph as compared to MRTG.

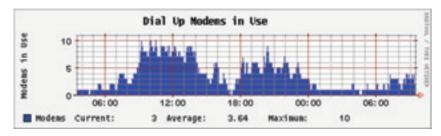


Figure 6.15: RRDtool gives you a lot of flexibility in how your collected network data may be displayed.

RRDtool is included in virtually all modern Linux distributions, and can be downloaded from *http://oss.oetiker.ch/rrdtool/.* 

#### ntop

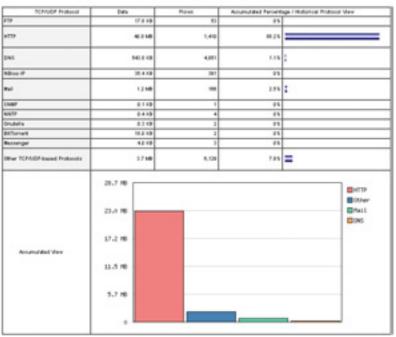
*http://www.ntop.org/*. For historical traffic analysis and usage, you will certainly want to investigate **ntop**. This program builds a detailed real-time report on observed network traffic, displayed in your web browser. It integrates with rrdtool, and makes graphs and charts visually depicting how the network is being used. On very busy networks, ntop can use a lot of CPU and disk space, but it gives you extensive insight into how your network is being used. It runs on Linux, BSD, Mac OS X, and Windows.

Some of its more useful features include:

- Traffic display can be sorted by various criteria (source, destination, protocol, MAC address, etc.).
- · Traffic statistics grouped by protocol and port number
- · An IP traffic matrix which shows connections between machines
- Network flows for routers or switches that support the NetFlow protocol
- Host operating system identification

- P2P traffic identification
- · Numerous graphical charts
- Perl, PHP, and Python API

Ntop is available from *http://www.ntop.org/* and is available for most operating systems. It is often included in many of the popular Linux distributions, including RedHat, Debian, and Ubuntu. While it can be left running to collect historical data, ntop can be fairly CPU intensive, depending on the amount of traffic observed. If you are going to run it for long periods you should monitor the CPU utilization of the monitoring machine.



**Global TCP UDP Protocol Distribution** 

Figure 6.16: ntop displays a wealth of information about how your network is utilized by various clients and servers.

The main disadvantage of ntop is that it does not provide instantaneous information, only long-term totals and averages. This can make it difficult to use to diagnose a problem that starts suddenly.

#### Cacti

*http://www.cacti.net/.* **Cacti** is a front-end for RRDtool. It stores all of the necessary information to create graphs in a MySQL database. The front-end is written in PHP. Cacti does the work of maintaining graphs, data sources,

and handles the actual data gathering. There is support for SNMP devices, and custom scripts can easily be written to poll virtually any conceivable network event.

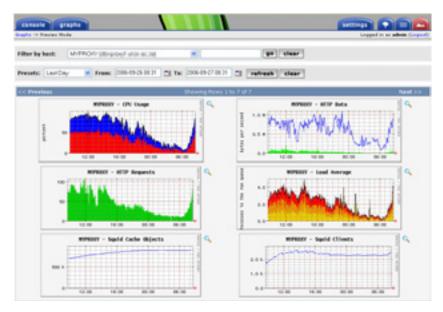


Figure 6.17: Cacti can manage the polling of your network devices, and can build very complex and informative visualizations of network behavior.

Cacti can be somewhat confusing to configure, but once you work through the documentation and examples, it can yield very impressive graphs. There are hundreds of templates for various systems available on the cacti website, and the code is under rapid development.

#### NetFlow

NetFlow is a protocol for collecting IP traffic information invented by Cisco. From the Cisco website:

Cisco IOS NetFlow efficiently provides a key set of services for IP applications, including network traffic accounting, usage-based network billing, network planning, security, Denial of Service monitoring capabilities, and network monitoring. NetFlow provides valuable information about network users and applications, peak usage times, and traffic routing.

Cisco routers can generate NetFlow information which is available from the router in the form of UDP packets. NetFlow is also less CPU-intensive on Cisco routers than using SNMP. It also provides more granular information than SNMP, letting you get a more detailed picture of port and protocol usage.

This information is collected by a NetFlow collector that stores and presents the data as an aggregate over time. By analyzing flow data, one can build a picture of traffic flow and traffic volume in a network or on a connection. There are several commercial and free NetFlow collectors available. Ntop is one free tool that can act as a NetFlow collector and probe. Another is Flowc (see below).

It can also be desirable to use Netflow as a spot check tool, by just looking at a quick snapshot of data during a network crisis. Think of NetFlow as an alternative to SNMP for Cisco devices. For more information about NetFlow, see <a href="http://en.wikipedia.org/wiki/Netflow">http://en.wikipedia.org/wiki/Netflow</a>.

#### Flowc

http://netacad.kiev.ua/flowc/. Flowc is an open source NetFlow collector (see NetFlow above). It is lightweight and easy to configure. Flowc uses a MySQL database to store aggregated traffic information. Therefore, it is possible to generate your own reports from the data using SQL, or use the included report generators. The built-in report generators produce reports in HTML, plain text or a graphical format.

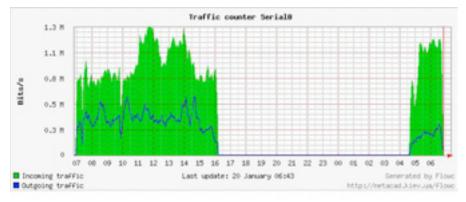


Figure 6.18: A typical flow chart generated by Flowc.

The large gap in data probably indicates a network outage. Trending tools typically will not notify you of outages, but merely log the occurrence. To be notified when network problems occur, use a realtime monitoring tool such as Nagios (see **Page 200**).

# SmokePing

http://oss.oetiker.ch/smokeping/. **SmokePing** is a deluxe latency measurement tool written in Perl. It can measure, store and display latency, latency distribution and packet loss all on a single graph. SmokePing uses RRDtool for data storage, and can draw very informative graphs that present up to the minute information on the state of your network connection. It is very useful to run SmokePing on a host with good connectivity to your entire network. Over time, trends are revealed that can point to all sorts of network problems. Combined with MRTG (see **Page 190**) or Cacti (see **Page 192**), you can observe the effect that network congestion has on packet loss and latency. SmokePing can optionally send alerts when certain conditions are met, such as when excessive packet loss is seen on a link for an extended period of time. An example of SmokePing in action is shown in **Figure 6.19**.

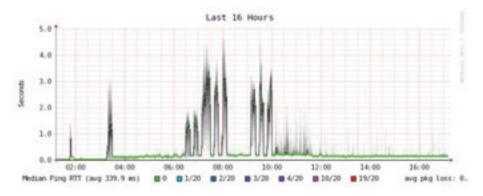


Figure 6.19: SmokePing can simultaneously display packet loss and latency spreads in a single graph.

# EtherApe

*http://etherape.sourceforge.net/.* **EtherApe** displays a graphical representation of network traffic. Hosts and links change size depending on the amount of traffic sent and received. The colors change to represent the protocol most used. As with wireshark and tcpdump, data can be captured "off the wire" from a live network connection or read from a tcpdump capture file.

EtherApe doesn't show quite as much detail as ntop, but its resource requirements are much lighter.

# iptraf

*http://iptraf.seul.org/.* **IPTraf** is a lightweight but powerful LAN monitor. It has an ncurses interface and runs in a command shell. IPTraf takes a moment to measure observed traffic, and then displays various network statistics including TCP and UDP connections, ICMP and OSPF information, traffic flows, IP checksum errors, and more. It is a simple to use program that uses minimal system resources.

While it does not keep historical data, it is very useful for displaying an instantaneous usage report.

IPTraf						
Proto/Port	🗕 Pkts 🗕	🗕 Bytes 🗕	– PktsTo 🗕	BytesTo	PktsFrom By	/tesFrom 🗕
TCP/80	23	12534	10	559	13	11975
UDP/137	22	1716	11	858	11	858
UDP/53	104	14635	61	4591	43	10044
TCP/25	460	78861	247	52772	213	25289
TCP/53	4	240	4	240	0	0
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U0P/138	12	2762	6	1381	6	1381
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Figure 6.20: iptraf's statistical breakdown of traffic by port.

#### Argus

*http://qosient.com/argus/.* **Argus** stands for **Audit Record Generation and Utilization System**. Argus is also the name of the mythological Greek god who had hundreds of eyes.

From the Argus website:

Argus generates flow statistics such as connectivity, capacity, demand, loss, delay, and jitter on a per transaction basis. Argus can be used to analyze and report on the contents of packet capture files or it can run as a continuous monitor, examining data from a live interface; generating an audit log of all the network activity seen in the packet stream. Argus can be deployed to monitor individual end-systems, or an entire enterprises network activity. As a continuous monitor, Argus provides both push and pull data handling models, to allow flexible strategies for collecting network audit data. Argus data clients support a range of operations, such as sorting, aggregation, archival and reporting.

Argus consists of two parts: a master collector that reads packets from a network device, and a client that connects to the master and displays the usage statistics. Argus runs on BSD, Linux, and most other UNIX systems.

# NeTraMet

http://freshmeat.net/projects/netramet/. **NeTraMet** is another popular flow analysis tool. Like Argus, NeTraMet consists of two parts: a collector that gathers statistics via SNMP, and a manager that specifies which flows should be watched. Flows are specified using a simple programming language that define the addresses used on either end, and can include Ethernet, IP, protocol information, or other identifiers. NeTraMet runs on DOS and most UNIX systems, including Linux and BSD.

# Throughput testing

How fast can the network go? What is the actual usable capacity of a particular network link? You can get a very good estimate of your throughput capacity by flooding the link with traffic and measuring how long it takes to transfer the data.



Figure 6.21: Tools such as this one from SpeedTest.net are pretty, but don't always give you an accurate picture of network performance.

While there are web pages available that will perform a "speed test" in your browser (such as *http://www.dslreports.com/stest* or *http://speedtest.net/*), these tests are increasingly inaccurate as you get further from the testing source. Even worse, they do not allow you to test the speed of a given link, but only the speed of your link to a particular site on the Internet. Here are a few tools that will allow you to perform throughput testing on your own networks.

#### ttcp

*http://ftp.arl.mil/ftp/pub/ttcp/*. Now a standard part of most Unix-like systems, **ttcp** is a simple network performance testing tool. One instance is run on either side of the link you want to test. The first node runs in receive mode, and the other transmits:

node\_a\$ ttcp -r -s
node\_b\$ ttcp -t -s node\_a
ttcp-t: buflen=8192, nbuf=2048, align=16384/0, port=5001 tcp -> node\_a
ttcp-t: socket
ttcp-t: connect
ttcp-t: 16777216 bytes in 249.14 real seconds = 65.76 KB/sec +++
ttcp-t: 2048 I/O calls, msec/call = 124.57, calls/sec = 8.22
ttcp-t: 0.0user 0.2sys 4:09real 0% 0i+0d 0maxrss 0+0pf 7533+0csw

After collecting data in one direction, you should reverse the transmit and receive partners to test the link in the other direction. It can test UDP as well as TCP streams, and can alter various TCP parameters and buffer lengths to give the network a good workout. It can even use a user-supplied data stream instead of sending random data. Remember that the speed readout is in kilobytes, not kilobits. Multiply the result by 8 to find the speed in kilobits per second.

The only real disadvantage to ttcp is that it hasn't been developed in years. Fortunately, the code has been released in the public domain and is freely available. Like ping and traceroute, ttcp is found as a standard tool on many systems.

#### iperf

*http://dast.nlanr.net/Projects/Iperf/*. Much like ttcp, **iperf** is a commandline tool for estimating the throughput of a network connection. It supports many of the same features as ttcp, but uses a "client" and "server" model instead of a "receive" and "transmit" pair. To run iperf, launch a server on one side and a client on the other:

The server side will continue to listen and accept client connections on port 5001 until you hit control-C to kill it. This can make it handy when running multiple test runs from a variety of locations.

The biggest difference between ttcp and iperf is that iperf is under active development, and has many new features (including IPv6 support). This makes it a good choice as a performance tool when building new networks.

## bing

*http://fgouget.free.fr/bing/index-en.shtml*. Rather than flood a connection with data and see how long the transfer takes to complete, **Bing** attempts to estimate the available throughput of a point-to-point connection by analyzing round trip times for various sized ICMP packets. While it is not always as accurate as a flood test, it can provide a good estimate without transmitting a large number of bytes.

Since bing works using standard ICMP echo requests, so it can estimate available bandwidth without the need to run a special client on the other end, and can even attempt to estimate the throughput of links outside your network. Since it uses relatively little bandwidth, bing can give you a rough idea of network performance without running up the charges that a flood test would certainly incur.

# **Realtime tools**

It is desirable to find out when people are trying to break into your network, or when some part of the network has failed. Because no system administrator can be monitoring a network all the time, there are programs that constantly monitor the status of the network and can send alerts when notable events occur. The following are some open source tools that can help perform this task.

## Snort

**Snort** (*http://www.snort.org/*) is a packet sniffer and logger which can be used as a lightweight network intrusion detection system. It features rulebased logging and can perform protocol analysis, content searching, and packet matching. It can be used to detect a variety of attacks and probes, such as stealth port scans, CGI attacks, SMB probes, OS fingerprinting attempts, and many other kinds of anomalous traffic patterns. Snort has a realtime alert capability that can notify administrators about problems as they occur with a variety of methods.

Installing and running Snort is not trivial, and depending on the amount of network traffic, will likely require a dedicated monitoring machine with considerable resources. Fortunately, Snort is very well documented and has a strong user community. By implementing a comprehensive Snort rule set, you can identify unexpected behavior that would otherwise mysteriously eat up your Internet bandwidth. See *http://snort.org/docs/* for an extensive list of installation and configuration resources.

# Apache: mod\_security

ModSecurity (*http://www.modsecurity.org/*) is an open source intrusion detection and prevention engine for web applications. This kind of security tool is also known as a **web application firewall**. ModSecurity increases web application security by protecting web applications from known and unknown attacks. It can be used on its own, or as a module in the Apache web server (*http://www.apache.org/*).

There are several sources for updated mod\_security rules that help protect against the latest security exploits. One excellent resource is GotRoot, which maintains a huge and frequently updated repository of rules:

#### http://gotroot.com/tiki-index.php?page=mod\_security+rules

Web application security is important in defending against attacks on your web server, which could result in the theft of valuable or personal data, or in the server being used to launch attacks or send spam to other Internet users. As well as being damaging to the Internet as a whole, such intrusions can seriously reduce your available bandwidth.

# Nagios

**Nagios** (http://nagios.org/) is a program that monitors hosts and services on your network, notifying you immediately when problems arise. It can send notifications via email, SMS, or by running a script, and will send notifications to the relevant person or group depending on the nature of the problem. Nagios runs on Linux or BSD, and provides a web interface to show up-to-the-minute system status.

Nagios is extensible, and can monitor the status of virtually any network event. It performs checks by running small scripts at regular intervals, and checks the results against an expected response. This can yield much more sophisticated checks than a simple network probe. For example, ping (**page 185**) may tell you that a machine is up, and nmap may report that a TCP port responds to requests, but Nagios can actually retrieve a web page or make a database request, and verify that the response is not an error.

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Figure 6.22: Nagios keeps you informed the moment a network fault or service outage occurs.

Nagios can even notify you when bandwidth usage, packet loss, machine room temperature, or other network health indicator crosses a particular threshold. This can give you advance warning of network problems, often allowing you to respond to the problem before users have a chance to complain.

# Zabbix

**Zabbix** (http://www.zabbix.org/) is an open source realtime monitoring tool that is something of a hybrid between Cacti and Nagios. It uses a SQL database for data storage, has its own graph rendering package, and performs all of the functions you would expect from a modern realtime monitor (such as SNMP polling and instant notification of error conditions). Zabbix is released under the GNU General Public License.

# Other useful tools

There are thousands of free network monitoring tools that fill very specialized needs. Here are a few of our favorites that don't quite fit into the above categories.

## Driftnet and Etherpeg.

These tools decode graphical data (such as GIF and JPEG files) and display them as a collage. As mentioned earlier, tools such as these are of limited use in

troubleshooting problems, but are very valuable for demonstrating the insecurity of unencrypted protocols. **Etherpeg** is available from *http://www.etherpeg.org/*, and **Driftnet** can be downloaded at *http://www.ex-parrot.com/~chris/driftnet/*.



Figure 6.23: A web collage generated by Etherpeg.

#### ngrep

**Ngrep** provides most of GNU grep's pattern matching features, but applies them to network traffic. It currently recognizes IPv4 and IPv6, TCP, UDP, ICMP, IGMP, PPP, SLIP, FDDI, Token Ring, and much more. As it makes extensive use of regular expression matches, it is a tool suited to advanced users or those that have a good knowledge of regular expressions.

But you don't necessarily need to be a regex expert to be able to make basic use of ngrep. For example, to view all packets that contain the string GET (presumably HTTP requests), try this:

# ngrep -q GET

Pattern matches can be constrained further to match particular protocols, ports, or other criteria using BPF filters. This is the filter language used by common packet sniffing tools, such as tcpdump and snoop. To view GET or POST strings sent to destination port 80, use this command line:

# ngrep -q 'GET|POST' port 80

By using ngrep creatively, you can detect anything from virus activity to spam email. You can download ngrep at *http://ngrep.sourceforge.net/.* 

# What is normal?

If you are looking for a definitive answer as to what your traffic patterns **should** look like, you are going to be disappointed. There is no absolute right answer to this question, but given some work you can determine what is normal for your network. While every environment is different, some of the factors that can influence the appearance of your traffic patterns are:

- The capacity of your Internet connection
- The number of users that have access to the network
- The social policy (byte charging, quotas, honor system, etc.).
- The number, types, and level of services offered
- The health of the network (presence of viruses, excessive broadcasts, routing loops, open email relays, denial of service attacks, etc.).
- The competence of your computer users
- The location and configuration of control structures (firewalls, proxy servers, caches, and so on)

This is not a definitive list, but should give you an idea of how a wide range of factors can affect your bandwidth patterns. With this in mind, let's look at the topic of baselines.

# Establishing a baseline

Since every environment is different, you need to determine for yourself what your traffic patterns look like under normal situations. This is useful because it allows you to identify changes over time, either sudden or gradual. These changes may in turn indicate a problem, or a potential future problem, with your network.

For example, suppose that your network grinds to a halt, and you are not sure of the cause. Fortunately, you have decided to keep a graph of broadcasts as a percentage of the overall network traffic. If this graph shows a sudden increase in the amount of broadcast traffic, it may mean that your network has been infected with a virus. Without an idea of what is "normal" for your network (a baseline), you would not be able to see that the number of broadcasts had increased, only that it was relatively high, which may not indicate a problem. Baseline graphs and figures are also useful when analyzing the effects of changes made to the network. It is often very useful to experiment with such changes by trying different possible values. Knowing what the baseline looks like will show you whether your changes have improved matters, or made them worse.

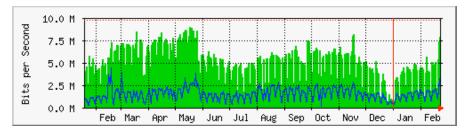


Figure 6.24: By collecting data over a long period of time, you can predict the growth of your network and make changes before problems develop.

In **Figure 6.24**, we can see the effect the implementation of delay pools has made on Internet utilization around the period of May. If we did not keep a graph of the line utilization, we would never know what the effect of the change over the long term was. When watching a total traffic graph after making changes, don't assume that just because the graph does not change radically that your efforts were wasted. You might have removed frivolous usage from your line only to have it replaced by genuine legitimate traffic. You could then combine this baseline with others, say the top 100 sites accessed or the average utilization by your top twenty users, to determine if habits have simply changed. As we will see later, MRTG, RRDtool, and Cacti are excellent tools you can use to keep a baseline.

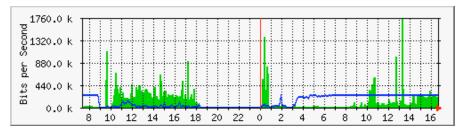


Figure 6.25: The traffic trend at Aidworld logged over a single day.

**Figure 6.25** shows traffic on an Aidworld firewall over a period of 24 hours. There is nothing apparently wrong with this graph, but users were complaining about slow Internet access.

**Figure 6.26** shows that the upload bandwidth use (dark area) was higher during working hours on the last day than on previous days. A period of heavy upload usage started every morning at 03:00, and was normally fin-

ished by 09:00, but on the last day it was still running at 16:30. Further investigation revealed a problem with the backup software, which ran at 03:00 every day.

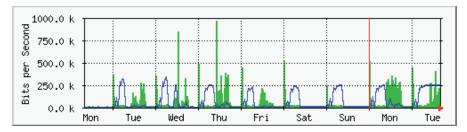


Figure 6.26: The same network logged over an entire week reveals a problem with backups, which caused unexpected congestion for network users.

**Figure 6.27** shows measurements of latency on the same connection as measured by a program called SmokePing. The position of the dots shows the average latency, while the gray smoke indicates the distribution of latency (jitter). The color of the dots indicates the number of lost packets. This graph over a period of four hours does not help to identify whether there are any problems on the network.

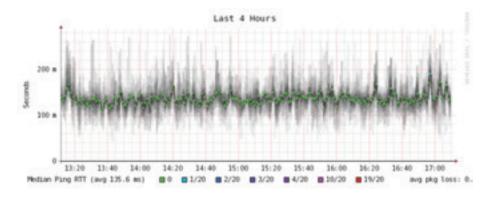


Figure 6.27: Four hours of jitter and packet loss.

The next graph (**Figure 6.28**) shows the same data over a period of 16 hours. This indicates that the values in the graph above are close to the normal level (baseline), but that there were significant increases in latency at several times during the early morning, up to 30 times the baseline value. This indicates that additional monitoring should be performed during these early morning periods to establish the cause of the high latency, which is probably heavy traffic of some kind.

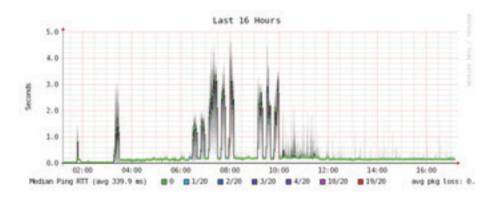


Figure 6.28: A higher spread of jitter is revealed in the 16 hour log.

**Figure 6.29** shows that Tuesday was significantly worse than Sunday or Monday for latency, especially during the early morning period. This might indicate that something has changed on the network.

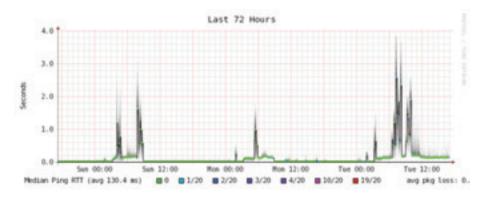


Figure 6.29: Zooming out to the week long view reveals a definite repetition of increased latency and packet loss in the early morning hours.

# How do I interpret the traffic graph?

In a basic network flow graph (such as that generated by the network monitor MRTG), the green area indicates *inbound traffic*, while the blue line indicates *outbound traffic*. Inbound traffic is traffic that originates from another network (typically the Internet) and is addressed to a computer inside your network. Outbound traffic is traffic that originates from your network, and is addressed to a computer somewhere on the Internet. Depending on what sort of network environment you have, the graph will help you understand how your network is actually being used. For example, monitoring of servers usually reveals larger amounts of outbound traffic as the servers respond to requests (such as sending mail or serving web pages), while monitoring cli-

ent machines might reveal higher amounts of inbound traffic to the machines as they receive data from the servers.

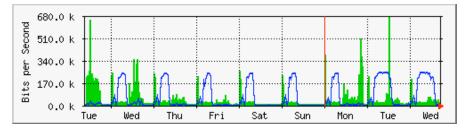


Figure 6.30: The classic network flow graph. The dark area represents inbound traffic, while the line represents outbound traffic. The repeating arcs of outbound traffic show when the nightly backups have run.

Traffic patterns will vary with what you are monitoring. A router will normally show more incoming traffic than outgoing traffic as users download data from the Internet. An excess of outbound bandwidth that is not transmitted by your network servers may indicate a peer-to-peer client, unauthorized server, or even a virus on one or more of your clients. There are no set metrics that indicate what outgoing traffic to incoming traffic should look like. It is up to you to establish a baseline to understand what normal network traffic patterns look like on your network.

# Detecting network overload

Figure 6.31 shows traffic on an overloaded Internet connection.

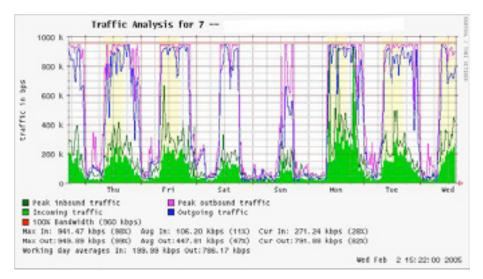


Figure 6.31: Flat-topped graphs indicate that a line is using the maximum available bandwidth, and is overutilized during these times.

The most apparent sign of overloading is the flat tops on outbound traffic during the middle of every day. Flat tops may indicate overloading, even if they are well below the maximum theoretical capacity of the link. In this case it may indicate that you are not getting as much bandwidth from your service provider as you expect.

## Measuring 95th percentile

The 95th percentile is a widely used mathematical calculation to evaluate regular and sustained utilization of a network pipe. Its value shows the highest consumption of traffic for a given period. Calculating the 95th percentile means that 95% of the time the usage is below a certain amount, and 5% of the time usage is above that amount. The 95th percentile is a good value to use to show the bandwidth that is actually used at least 95% of the time.

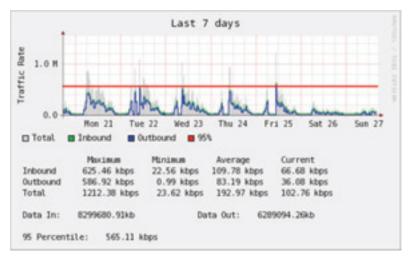


Figure 6.32: The horizontal line shows the 95th percentile amount.

MRTG and Cacti will calculate the 95th Percentile for you. This is a sample graph of a 960 kbps connection. The 95th percentile came to 945 kbps after discarding the highest 5% of traffic.

# Monitoring RAM and CPU usage

By definition, servers provide critical services that should always be available. Servers receive and respond to client machine requests, providing access to services that are the whole point of having a network in the first place. Therefore, servers must have sufficient hardware capabilities to accommodate the work load. This means they must have adequate RAM, storage, and processing power to accommodate the number of client requests. Otherwise, the server will take longer to respond, or in the worst case, may be incapable of responding at all. Since hardware resources are finite, it is important to keep track of how system resources are being used. If a core server (such as a proxy server or email server) is overwhelmed by requests, access times become slow. This is often perceived by users as a network problem.

There are several programs that can be used to monitor resources on a server. The simplest method on a Windows machine is to access the Task Manager using the **Ctrl Alt + Del** keys, and then click on the Performance tab. On a Linux or BSD box, you can type **top** in a terminal window. To keep historical logs of such performance, MRTG or RRDtool (on **Page 190**) can also be used.

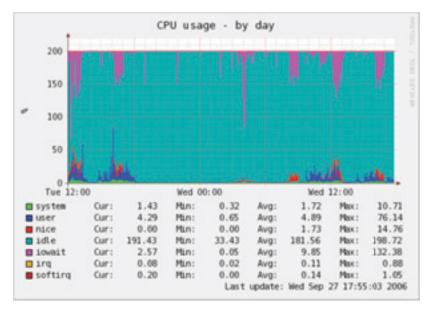


Figure 6.33: RRDtool can show arbitrary data, such as memory and CPU usage, expressed as an average over time.

Mail servers require adequate space, as some people may prefer to leave their email messages on the server for long periods of time. The messages can accumulate and fill the hard disk, especially if quotas are not in use. If the disk or partition used for mail storage fills up, the mail server cannot receive mail. If that disk is also used by the system, all kinds of system problems may occur as the operating system runs out of swap space and temporary storage.

File servers need to be monitored, even if they have large disks. Users will find a way to fill any size disk more quickly than you might think. Disk usage can be enforced through the use of quotas, or by simply monitoring usage and telling people when they are using too much. Nagios (see **Page 200**) can notify you when disk usage, CPU utilization, or other system resources cross a critical threshold.

If a machine becomes unresponsive or slow, and measurements show that a system resource is being heavily used, this may be an indication that an upgrade is required. If processor usage constantly exceeds 60% of the total, it may be time to upgrade the processor. Slow speeds could also be as a result of insufficient RAM. Be sure to check the overall usage of CPU, RAM, and disk space before deciding to upgrade a particular component.

A simple way to check whether a machine has insufficient RAM is to look at the hard disk light. When the light is on constantly, it usually means that the machine is constantly swapping large amounts of data to and from the disk. This is known as *thrashing*, and is extremely bad for performance. It can usually be fixed by investigating which process is using the most RAM, and killing or reconfiguring that process. Failing that, the system needs more RAM.

You should always determine whether it is more cost effective to upgrade an individual component or purchase a whole new machine. Some computers are difficult or impossible to upgrade, and it often costs more to replace individual components than to replace the entire system. Since the availability of parts and systems varies widely around the world, be sure to weigh the cost of parts vs. whole systems, including shipping and taxes, when determining the cost of upgrading.